

Stochastic Models and Analysis for Resource Management in Server Farms

Thesis Oral

VARUN GUPTA

Advantages of server farm architecture



Data center pods

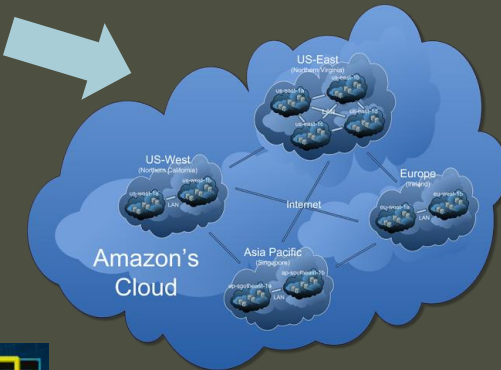


Supercomputers

- + high compute capacity
- + incremental growth
- + fault-tolerance
- + efficient resource utilization
- + energy efficiency
- + high parallelism

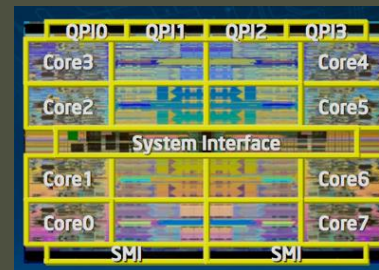


Array-of-Wimpy-Nodes



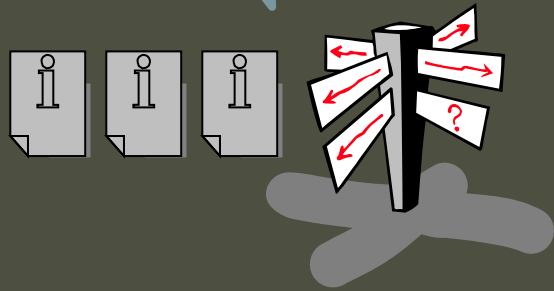
Amazon's Cloud

Cloud computing

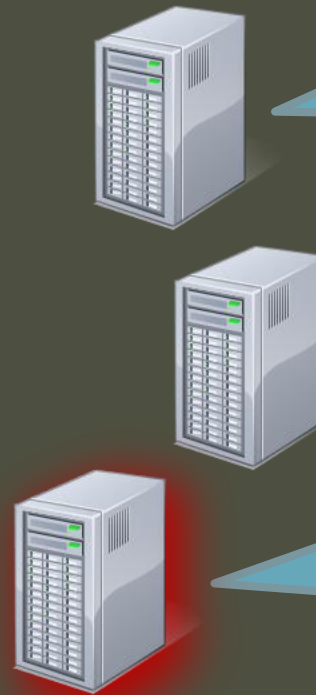


Multi-core chips

Design Choice 2:
Which server to assign jobs to?



Front-end load
balancer/dispatcher



Back-end
servers

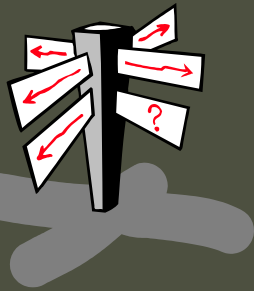
Design Choice 3:
Scheduling policy for
backend servers?

Design Choice 4:
When to turn servers
on/off for energy-
efficiency?

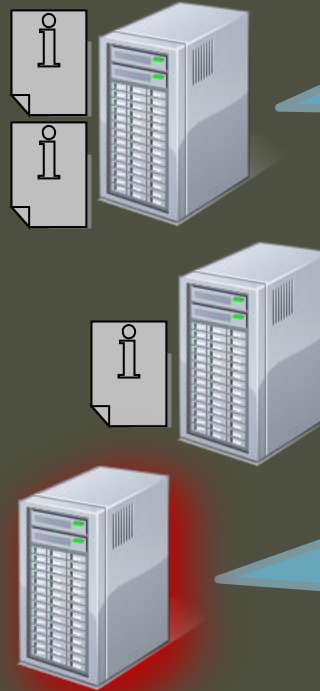
Design Choice 1: How many servers to buy? Of what capacity?



Design Choice 2:
Load Balancing policy



Front-end load
balancer/dispatcher



Back-end
servers

Design Choice 3:
Scheduling policy

Design Choice 4:
Dynamic capacity scaling

Design Choice 1: Provisioning/Dimensioning



Design Choice 2:
Load Balancing policy

Design Choice 3:
Scheduling policy

OPTION 1: Trial and error/Simulations

OPTION 2: Worst-case analysis

OPTION 3: Stochastic Modeling

- have estimates for real workloads
- understanding of “*what-if*” scenarios

Design Choice 4:
Dynamic capacity scaling

Design Choice 1: Provisioning/Dimensioning



Queueing Theory : The Origins



Manual telephone exchange (< 1900)

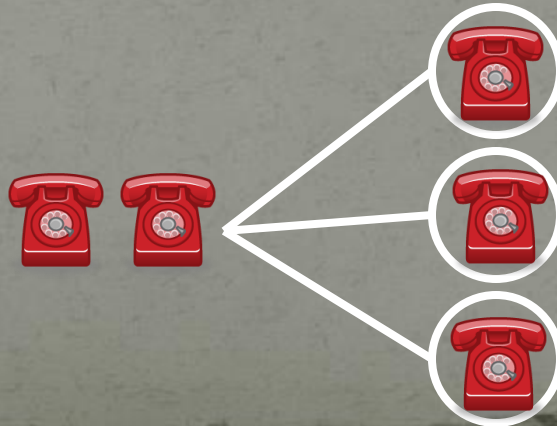


Automatic telephone exchange (~1910)



A.K. Erlang

Q: Use observed demand to dimension tel. exchanges



Queueing Theory : The Origins



Manual telephone exchange (< 1900)



Automatic telephone exchange (~1910)

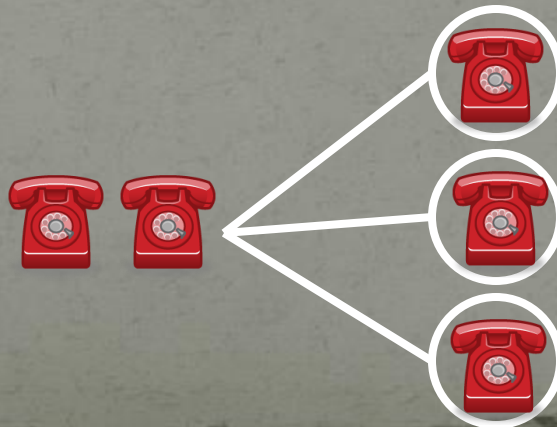


Congestion \Leftarrow stochastic demand



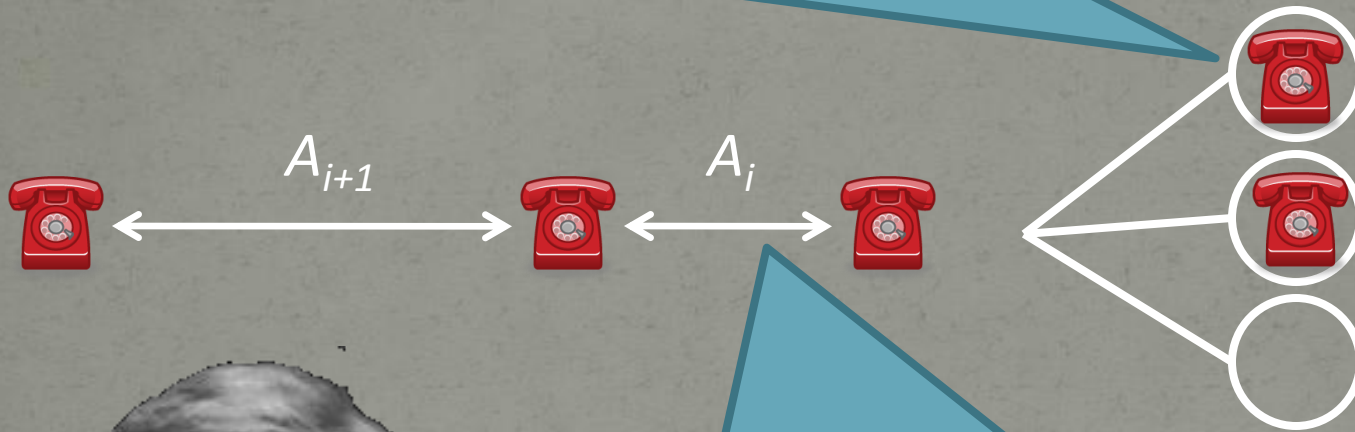
A.K. Erlang

Q: Use observed demand to dimension tel. exchanges



Queueing Theory : The Origins

Assumption 1: Call durations are i.i.d. Exponentially distributed random variables



A.K. Erlang

Assumption 2 (Poisson arrivals): Inter-call arrival times are i.i.d. Exponentially distributed

Q: Use observed demand to dimension tel. exchanges

Queueing Theory : The Origins

Assum

BUT existing queueing models are lacking for computing server farms

I. Workloads

- Classic models assume low variability in workload

II. Architectures

- Assume First-Come-First-Served servers
- Scale of traditional applications much smaller than data centers
- Dynamic capacity scaling not feasible

NEED new analysis and new models

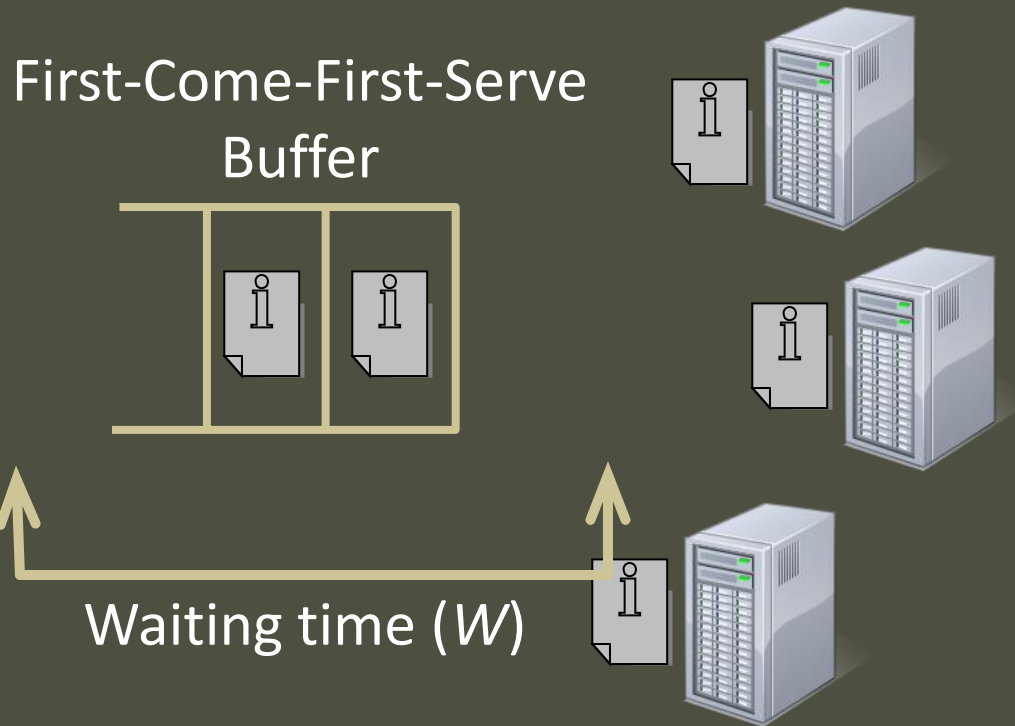
Part I. Impact of new workloads

- New analysis for a classical multi-server model
- Broader applications of analysis technique

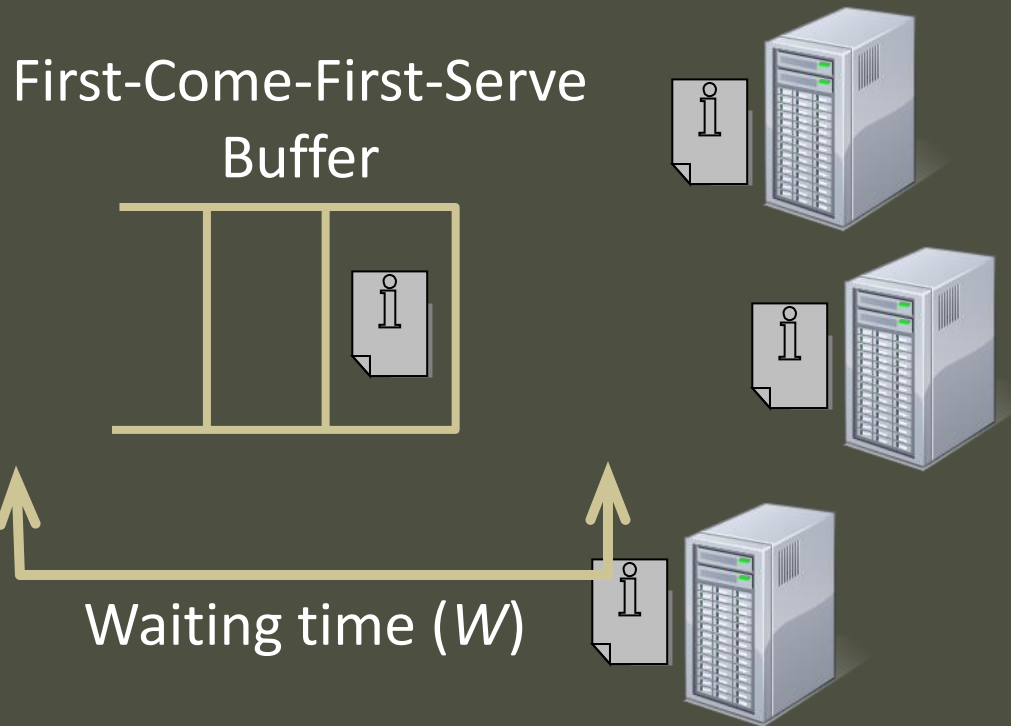
Part II. Impact of new architectures on:

- Concurrency control for servers
- Server management policies for energy-efficiency
- Load balancing

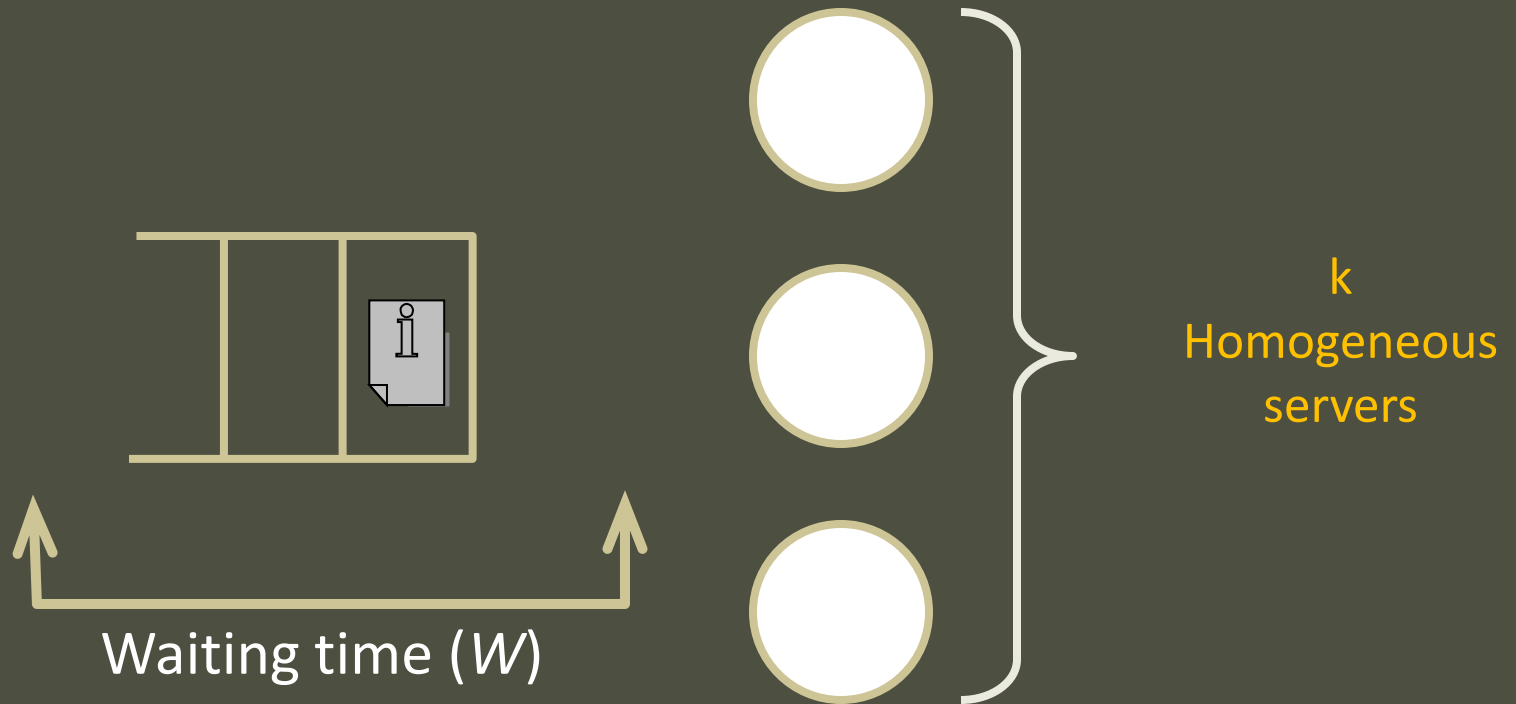
A classic multi-server model



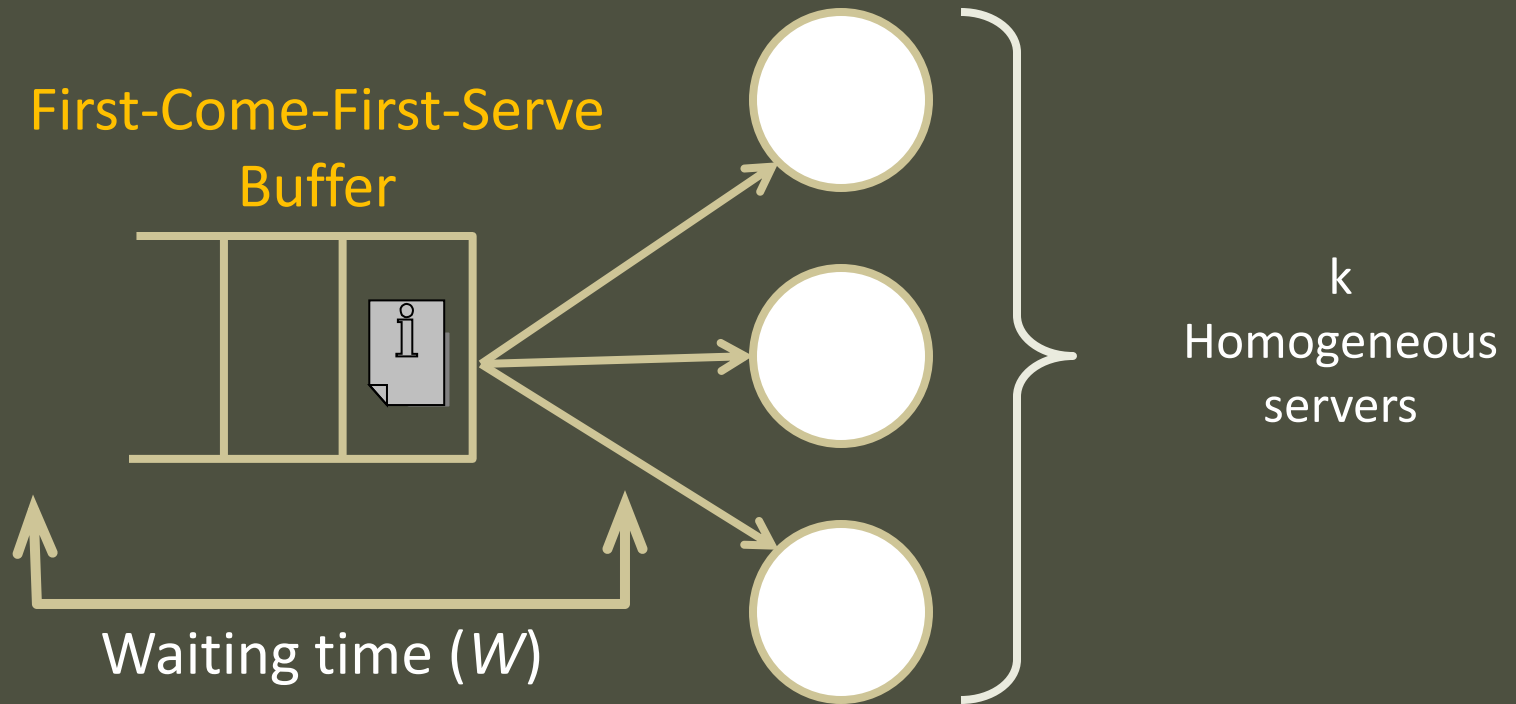
The $M/G/k/FCFS$ model



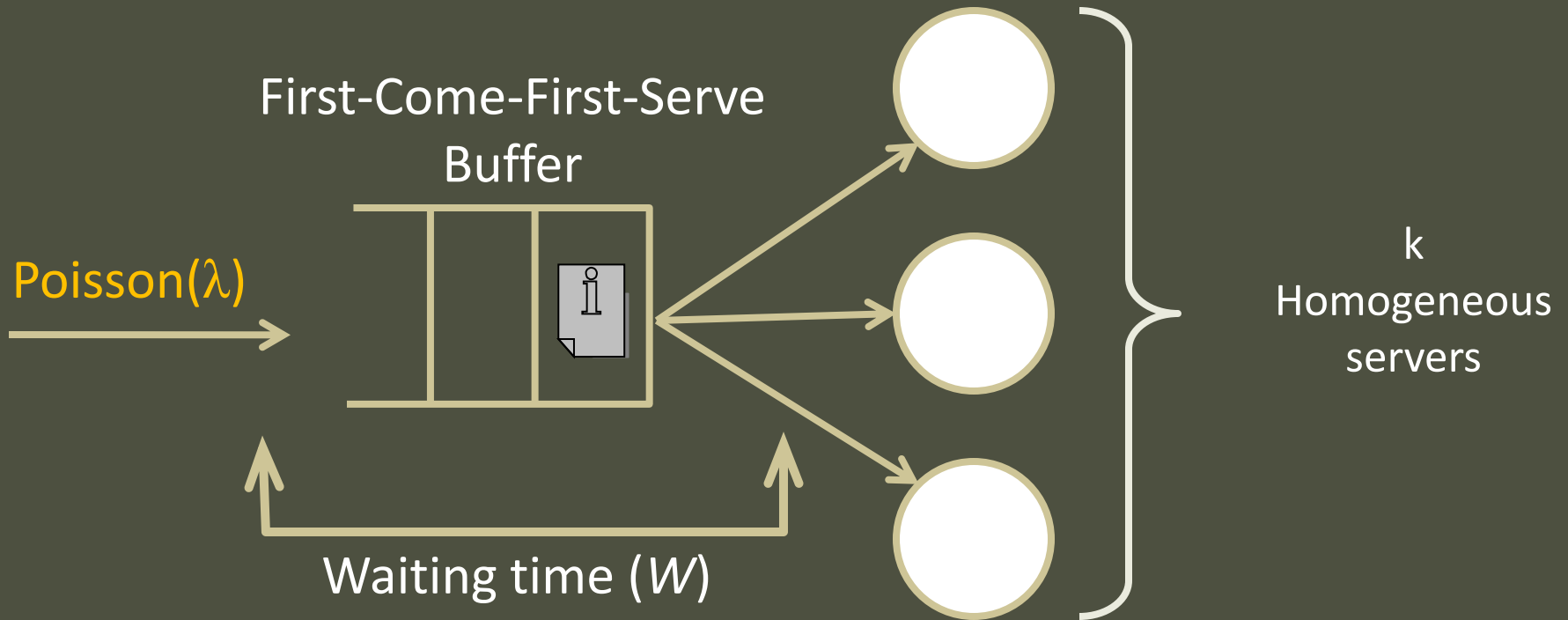
The $M/G/k/FCFS$ model



The $M/G/k/FCFS$ model

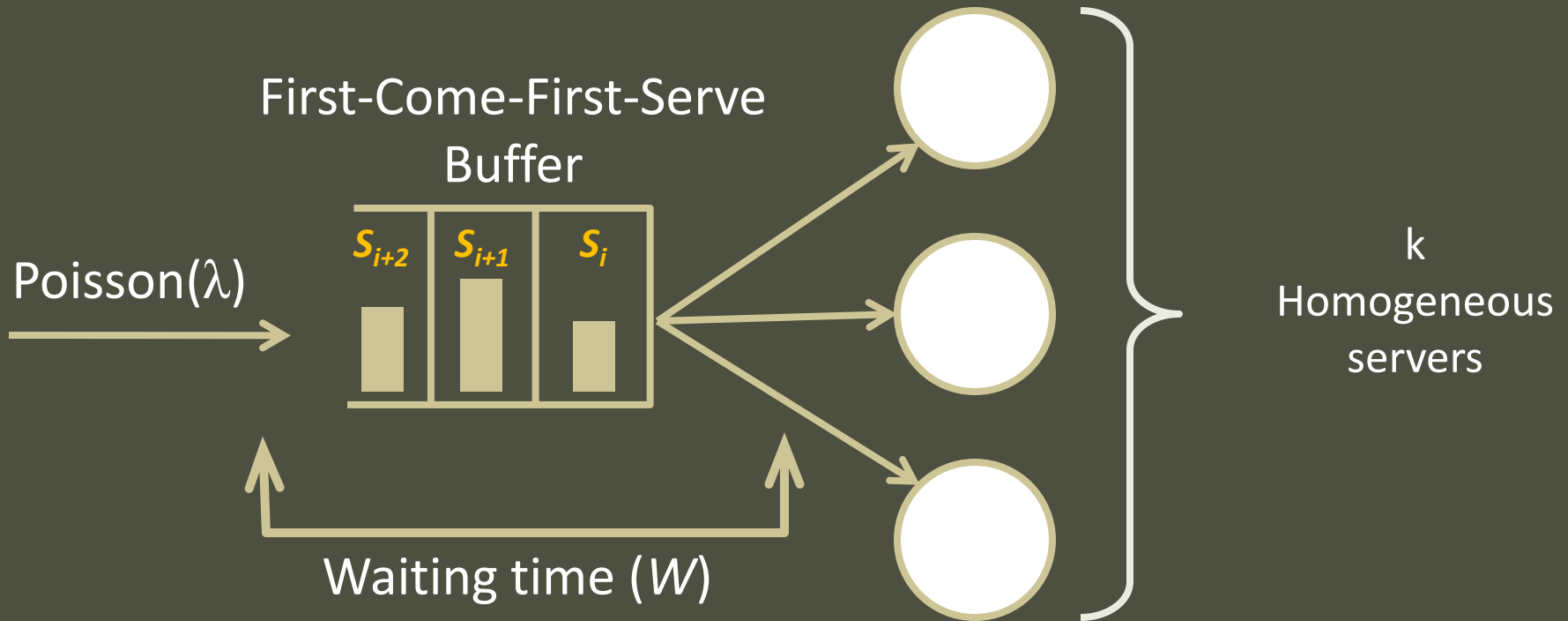


The $M/G/k/FCFS$ model



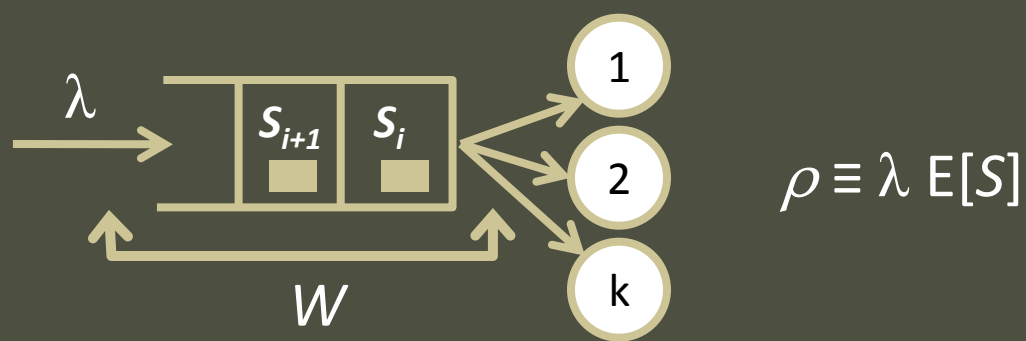
- λ = arrival rate

The $M/G/k/FCFS$ model



- λ = arrival rate
- job sizes (S_1, S_2, \dots) i.i.d. samples from S
- "load" $\rho \equiv \lambda E[S]$

GOAL : $E[W^{M/G/k}]$



k=1

Case : $S \sim$ Exponential ($M/M/1$)

Analyze $E[W^{M/M/1}]$ via Markov chain
(easy)

Case: $S \sim$ General ($M/G/1$)

$$E[W^{M/G/1}] = \frac{C^2+1}{2} E[W^{M/M/1}]$$

$$C^2 = \frac{\text{var}(S)}{E[S]^2}$$

Sq. Coeff. of Variation (SCV)
> 20 for computing workloads

k>1

Case : $S \sim$ Exponential ($M/M/k$)

$E[W^{M/M/k}]$ via Markov chain

Case: $S \sim$ General ($M/G/k$)

No exact analysis known

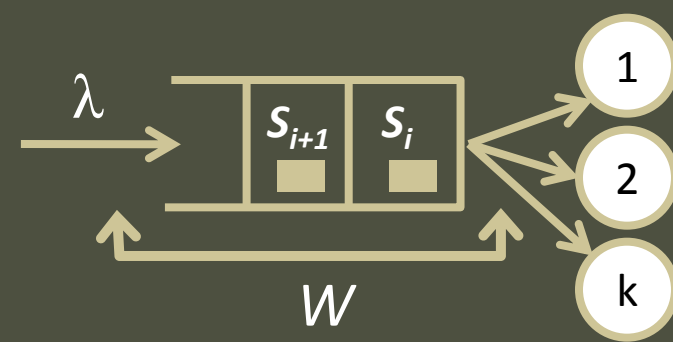
The Gold-standard approximation:

Lee, Longton (1959)

$$E[W^{M/G/k}] \approx \frac{C^2+1}{2} E[W^{M/M/k}]$$

Lee, Longton approximation:

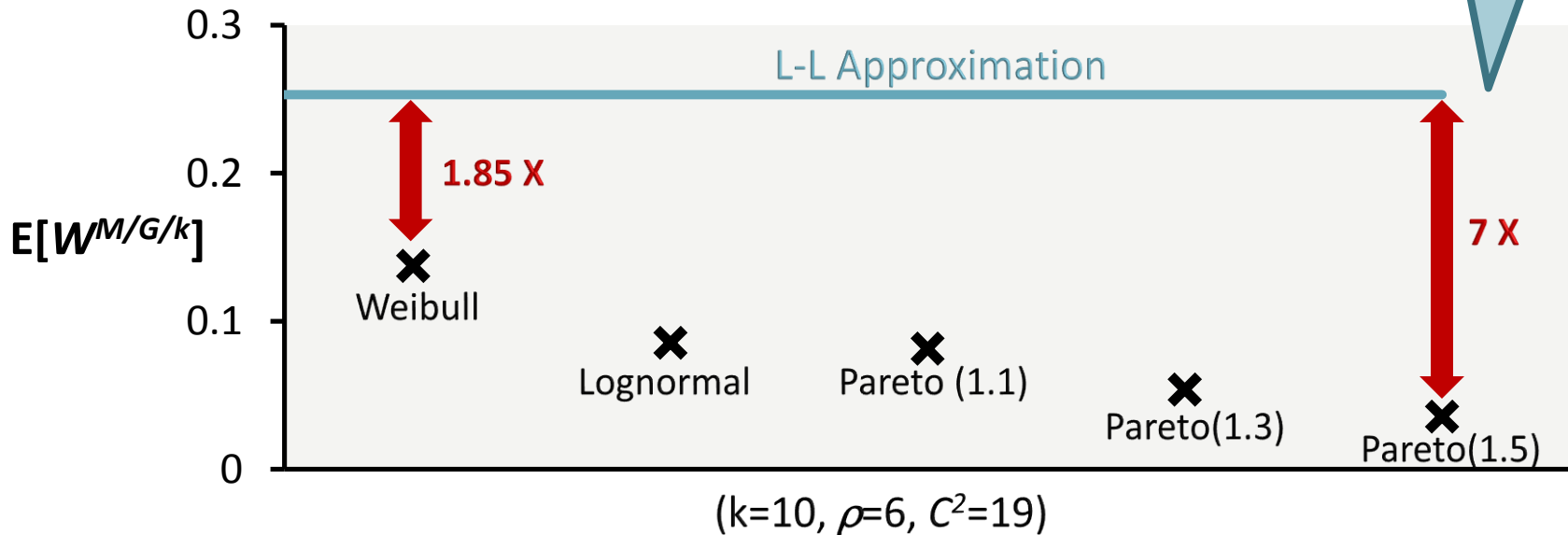
$$E[W^{M/G/k}] \approx \frac{C^2+1}{2} E[W^{M/M/k}]$$



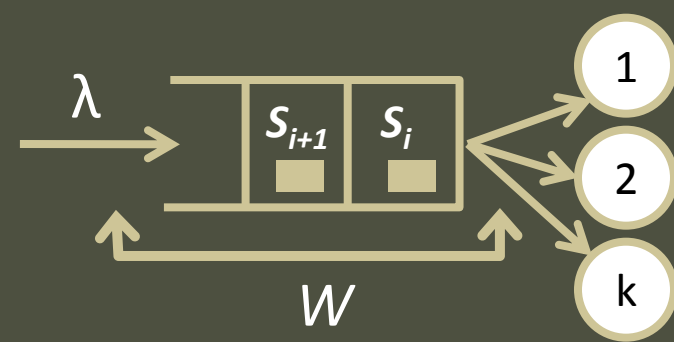
👍 Simple

👍 Exact for $k=1$

👍 Asymptotically tight as $\rho \rightarrow k$ (think Central Limit Thm.)



Outline: Part I

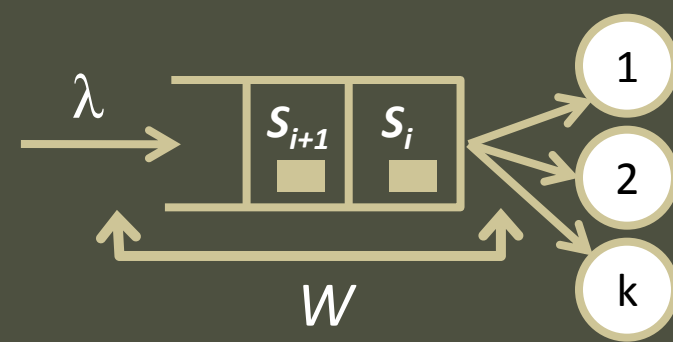


2 moments not enough for $E[W^{M/G/k}]$

Tighter bounds via higher moments of job size distribution

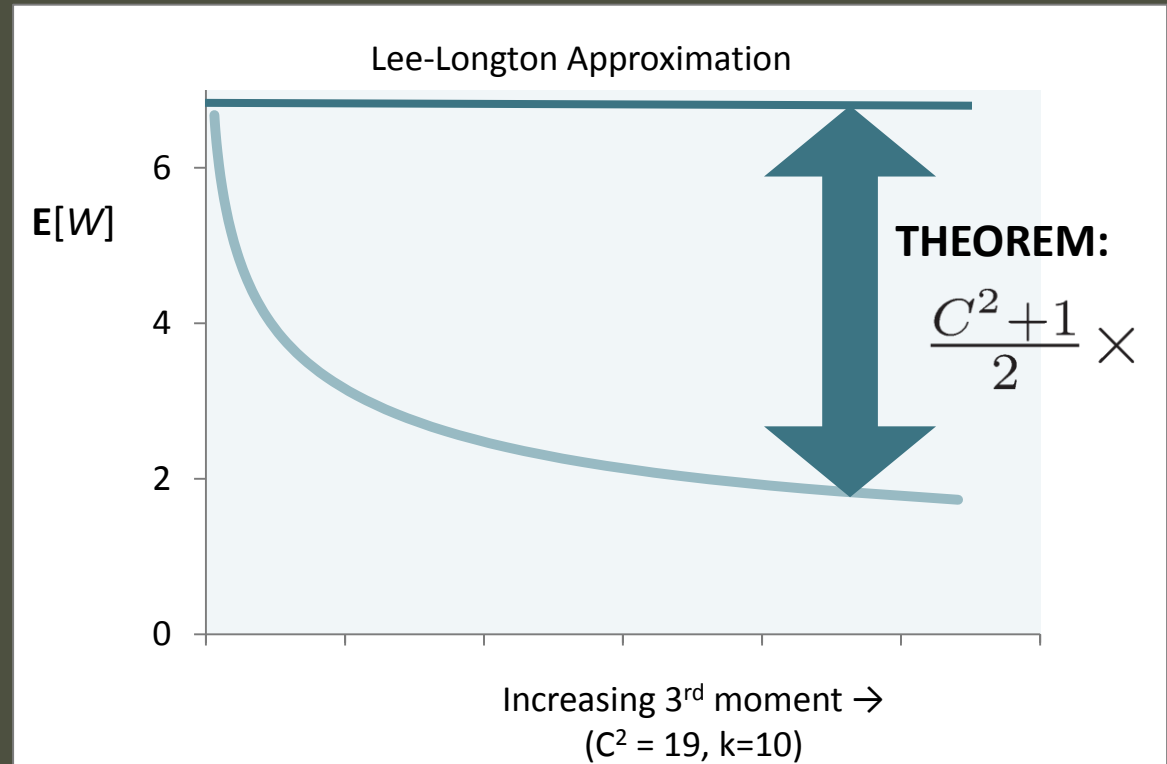
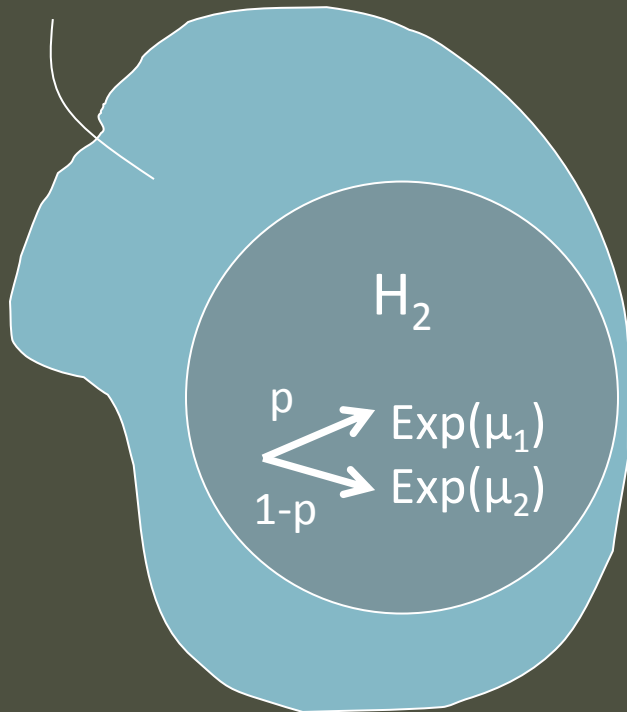
Lee, Longton approximation:

$$E[W^{M/G/k}] \approx \frac{C^2 + 1}{2} E[W^{M/M/k}]$$

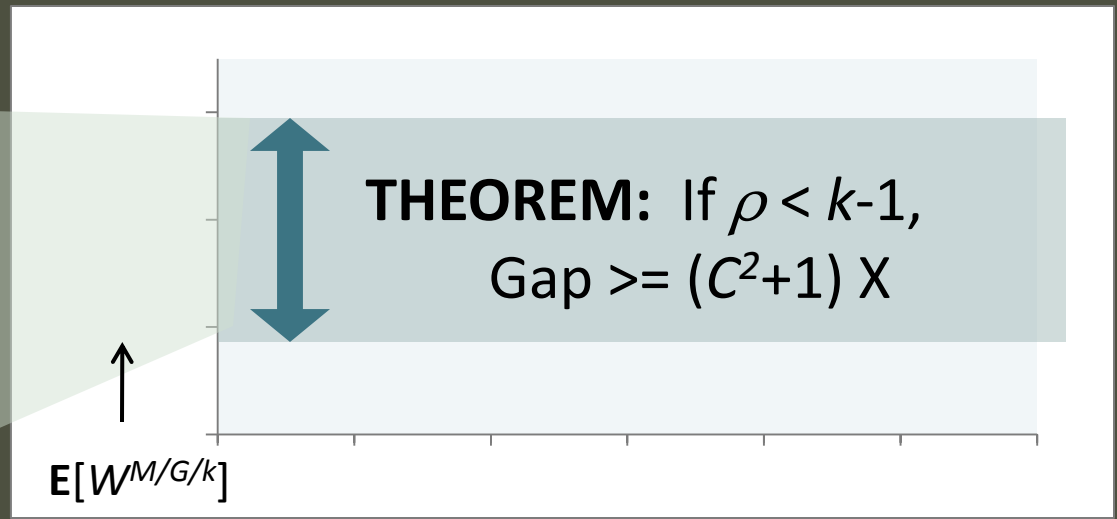


GOAL: Bounds on approximation ratio

{G | 2 moments}



{G | 2 moments}



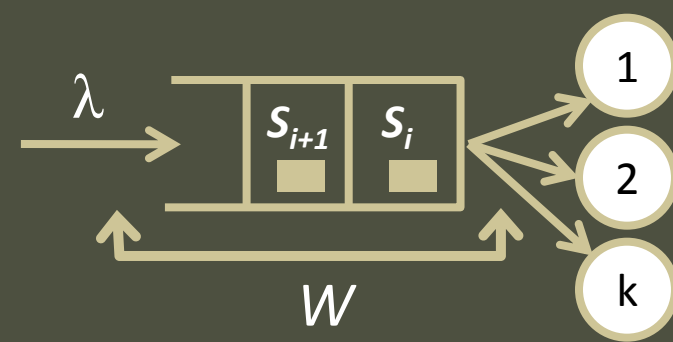
COR.: No approx. for $E[W^{M/G/k}]$ based on first two moments of job sizes can be accurate for all distributions when C^2 is large

PROOF: Analyze limit distributions in $D_2 \equiv$ mixture of 2 points



Approximations using higher moments?

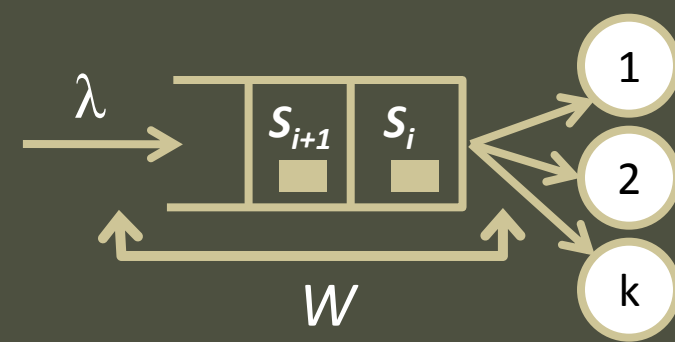
Outline: Part I



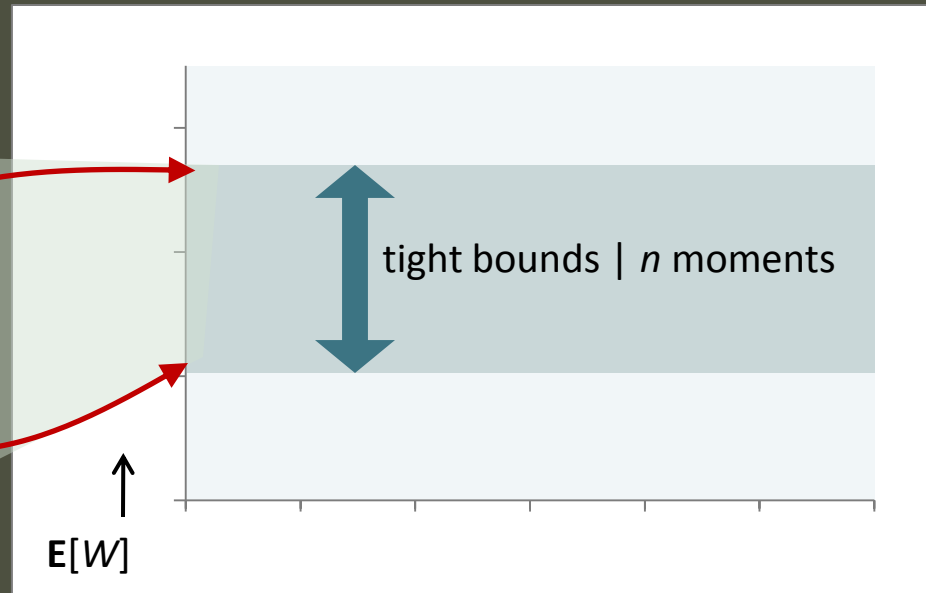
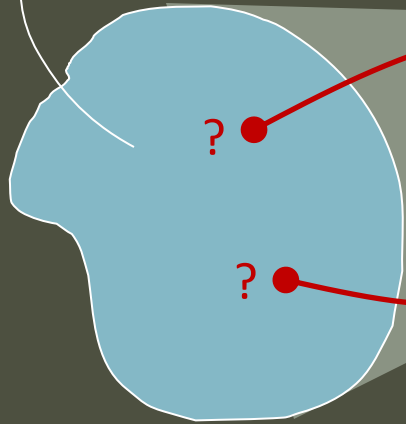
2 moments not enough for $E[W^{M/G/k}]$

Tighter bounds via higher moments of job size distribution

Exploiting higher moments



$\{G \mid n \text{ moments}\}$

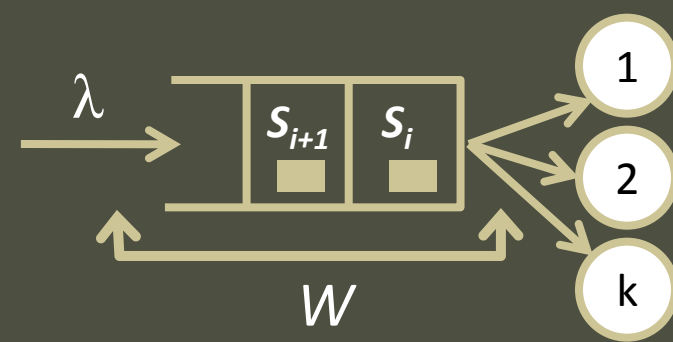


GOAL: Identify the “extremal” distributions with given moments

RELAXED GOAL: Extremal distributions in some “non-trivial” asymptotic regime

IDEA: Light-traffic asymptotics ($\lambda \rightarrow 0$)

Where we are...



GOAL: Tight bounds on $E[W^{M/G/k}]$ given n moments of S
IDEA: Identify extremal distributions

RELAXATION: Light Traffic

$$\lambda \rightarrow 0$$

Principal Representations and Extremal Problems

GIVEN: Moment conditions
on random variable X with
support $[0, B]$

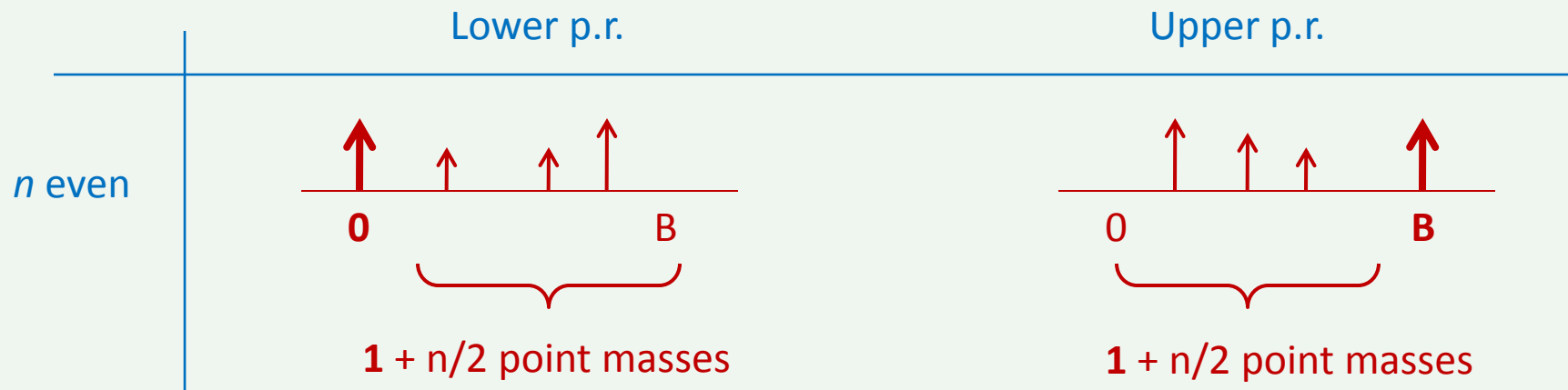
$$E[X^0] = m_0$$

$$E[X^1] = m_1$$

...

$$E[X^n] = m_n$$

Principal Representations (p.r.) on $[0, B]$ are distributions satisfying the moment conditions, and the following constraints on the support



Principal Representations and Extremal Problems

GIVEN: Moment conditions
on random variable X with
support $[0, B]$

$$E[X^0] = m_0$$

$$E[X^1] = m_1$$

...

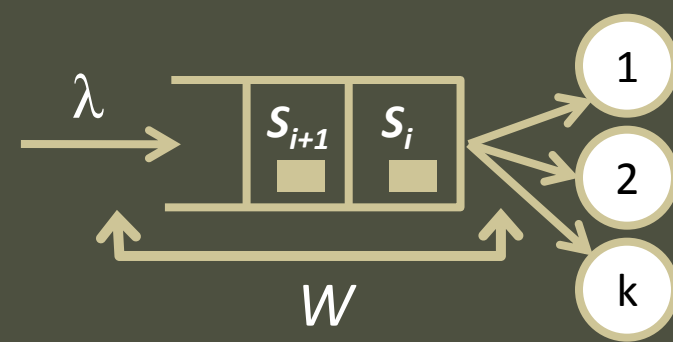
$$E[X^n] = m_n$$

Want to bound: $E[g(X)]$

THEOREM [Markov-Krein]:

If $\{x^0, \dots, x^n, g(x)\}$ is a Tchebycheff-system on $[0, B]$, then $E[g(X)]$ is extremized by the unique lower and upper principal representations of the moment sequence $\{m_0, \dots, m_n\}$.

Where we are...



GOAL: Tight bounds on $E[W^{M/G/k}]$ given n moments of S
IDEA: Identify extremal distributions

RELAXATION: Light Traffic

$$\lambda \rightarrow 0$$

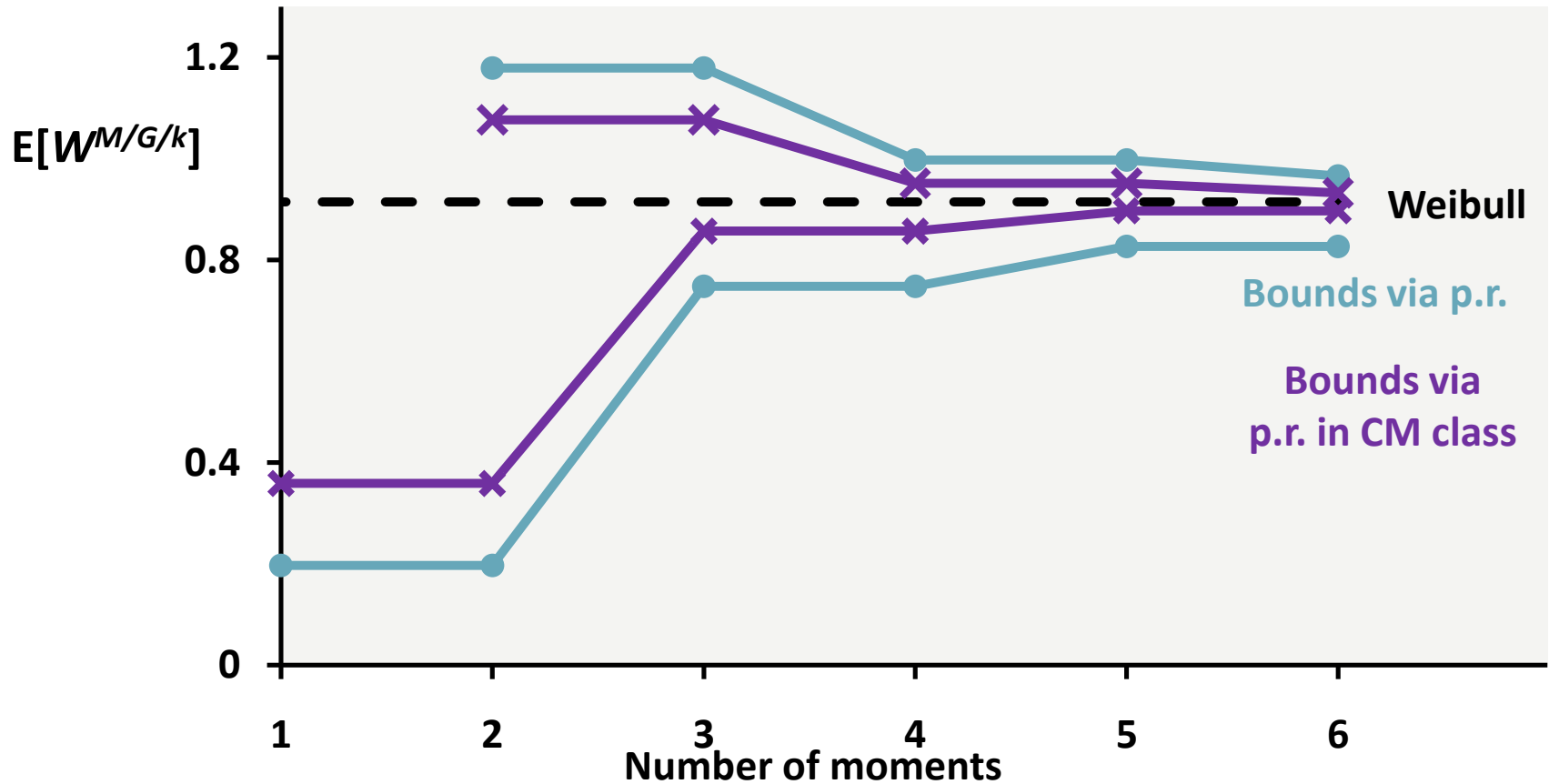
THEOREM:
For $n = 2$ or 3

RELAXATION 2: Restrict to Completely Monotone distributions (mixtures of Exponentials)

(contains Weibull, Pareto, Gamma)

THEOREM:
For all n .

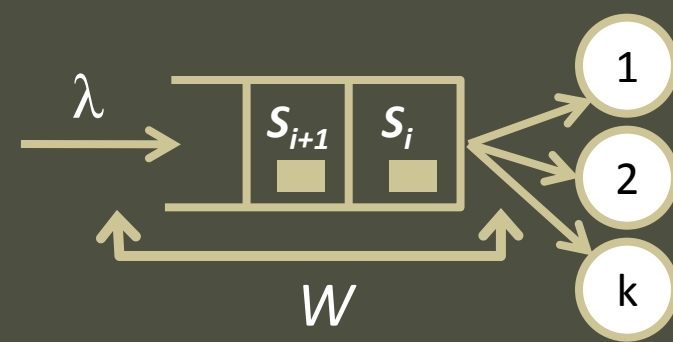
Simulation Results ($k=4, \rho=2.4,$)



Approximation Schema:

Refine **lower bound** via an additional **odd moment**,
Upper bound via **even moment** until gap is acceptable

Outline: Part I



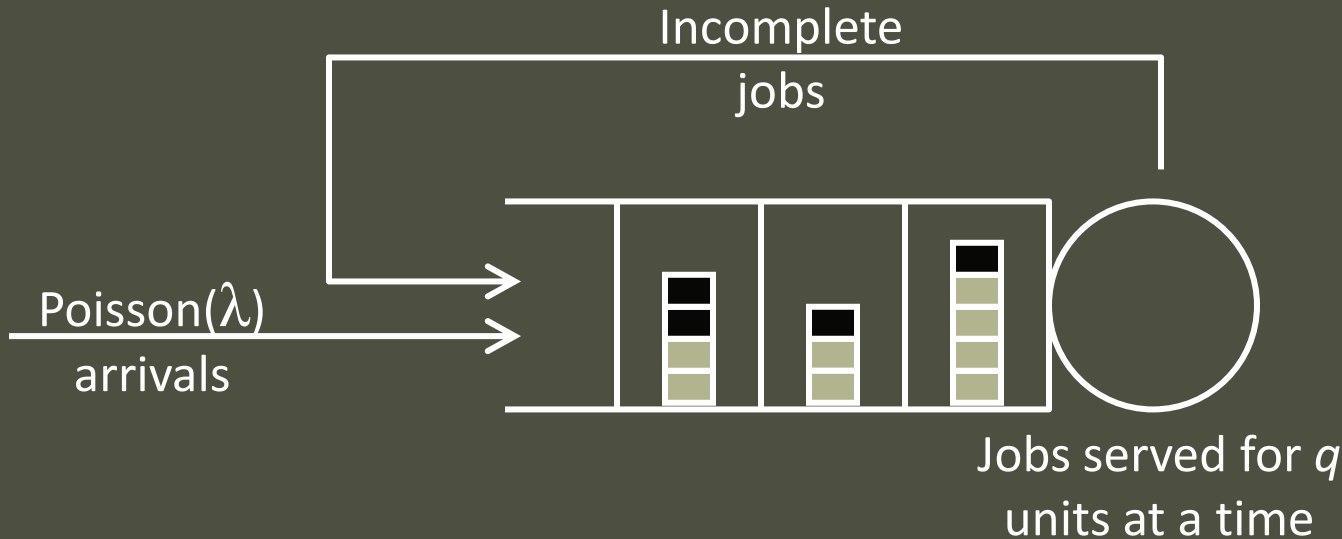
2 moments not enough for $E[W^{M/G/k}]$

Tighter bounds via higher moments of job size distribution

Many other “hard” queueing systems fit the approximation schema

Other queuing systems exhibiting Markov-Krein characterization

Example 1: M/G/1 Round-robin queue

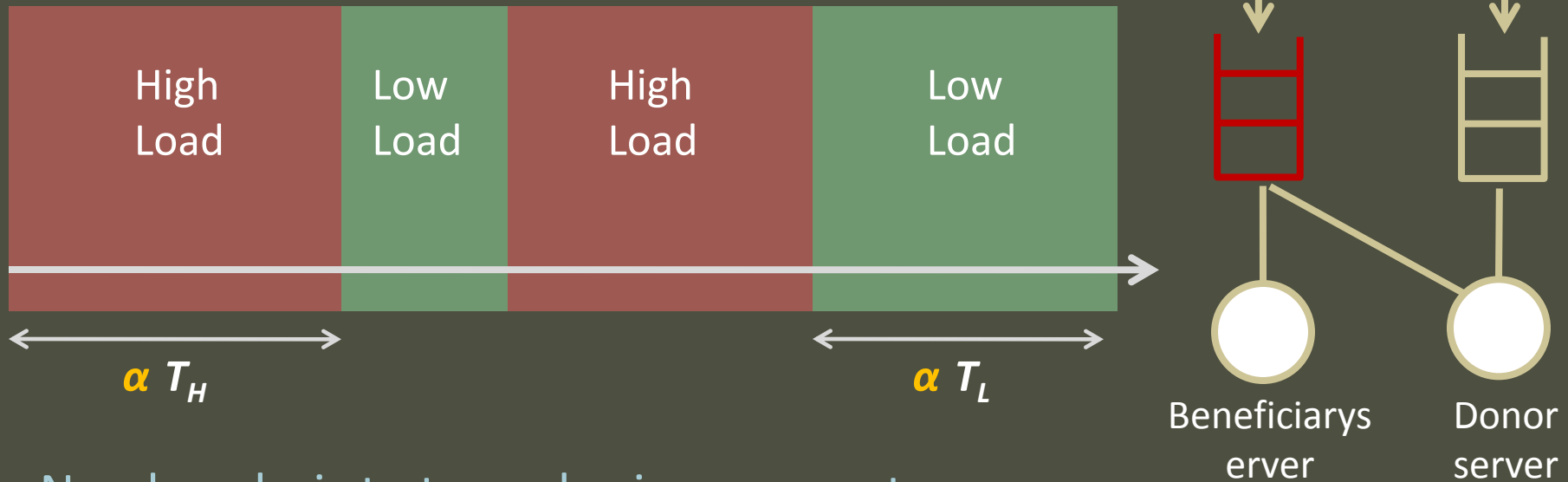


Need analysis to find q that balance overheads/performance

THEOREM: Upper and lower p.r. extremize mean response time under $\lambda \rightarrow 0$, when S is a mixture of Exponentials.

Other queuing systems exhibiting Markov-Krein characterization

Example 2: Systems with fluctuating load



Need analysis to tune sharing parameters

THEOREM: Upper and lower p.r. extremize mean waiting time under $\alpha \rightarrow 0$, when T_H, T_L are mixtures of Exponentials.

Part I. Impact of new workloads

- New analysis for a classical multi-server model
- Broader applications of analysis technique

Part II. Impact of new architectures on:

- Concurrency control for servers
- Dynamic server management for energy-efficiency
- Load balancing

TRADITIONAL

(e.g., manufacturing, call centers)

NEW

(Computing)

A.



FCFS servers



Processor sharing servers

B.

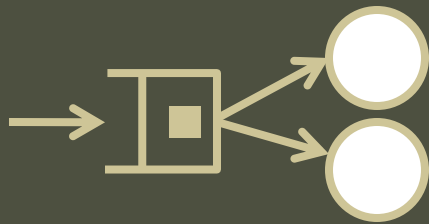


Ideal time-sharing

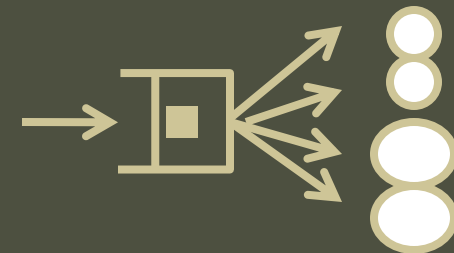


“Thrashing”

C.



Small + homogeneous farms



Large + heterogeneous farms

D.

Dynamic scaling for energy efficiency not feasible

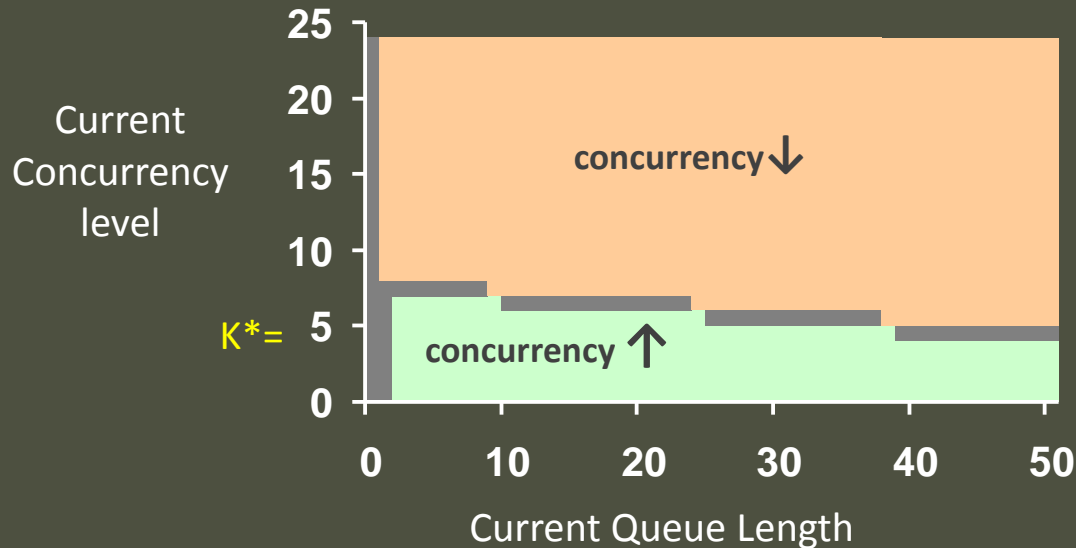
Servers with sleep states for energy efficiency

Application: Concurrency control in database servers

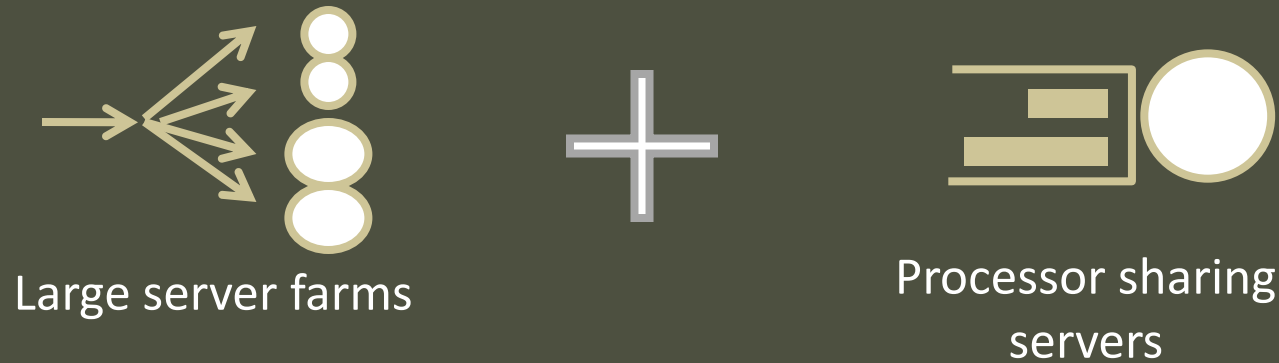


Contribution 1: Heuristic concurrency control algorithm under static arrival rate

Contribution 2: A simple traffic-oblivious heuristic



Application: Load Balancing in web server farms



Contribution 1: Join-the-Shortest-Queue (JSQ) near optimal for homogeneous servers

Contribution 2: JSQ is optimal for heterogeneous servers as size $\rightarrow \infty$

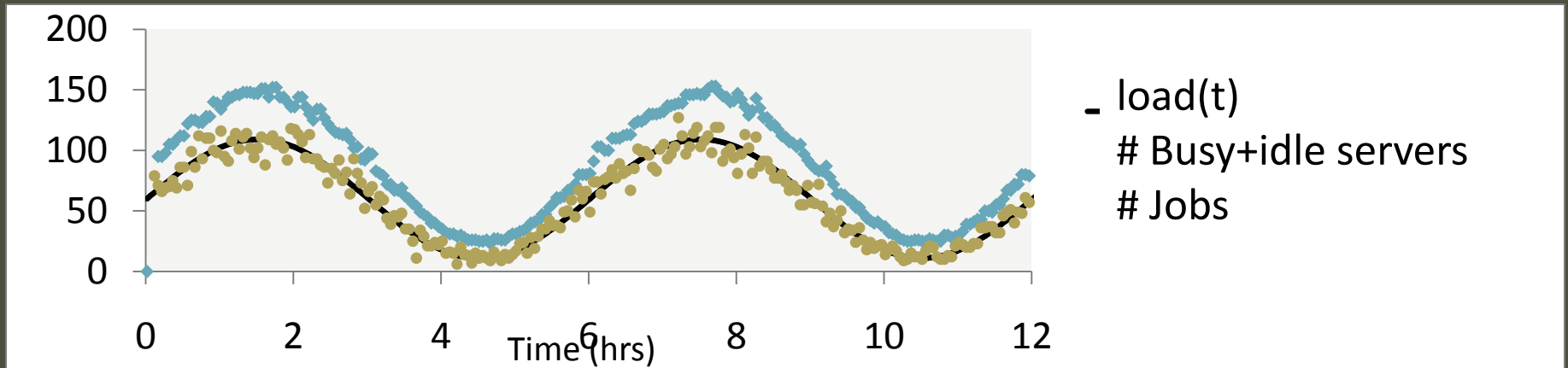
Contribution 3: First closed-form approximation for JSQ in many-servers regime

Application: Dynamic capacity scaling for energy-efficiency



No existing analysis for multi-server systems with setup delays

Contribution: A new traffic-oblivious policy **DELAYEDOFF**



DELAYEDOFF also extends to

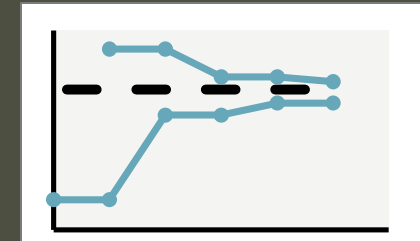
- Heterogeneous servers
- Virtual Machine management

Summary

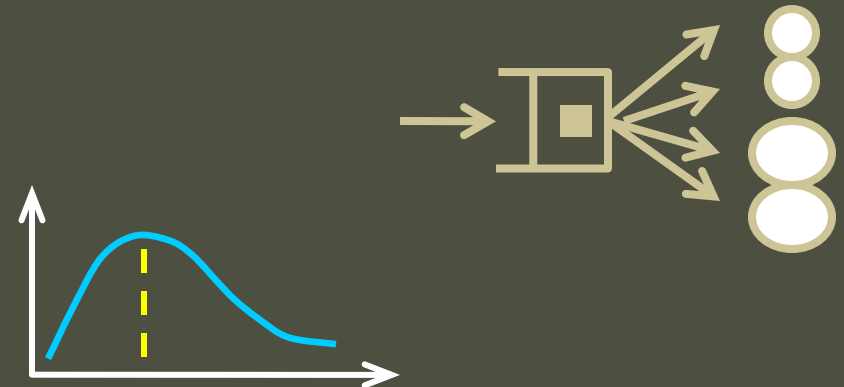
Stochastic modeling a powerful tool to analyze and optimize computer systems...

...but need new techniques to handle the new applications

• New workloads \Rightarrow new analysis



• New architectures \Rightarrow new models



References

[Performance'07] V. Gupta, M. Harchol-Balter, K. Sigman, and W. Whitt. Analysis of join-the-shortest-queue routing for web server farms. PERFORMANCE 2007

[Performance'10] V. Gupta, A. Gandhi, M. Harchol-Balter, and M. Kozuch. Optimality analysis of energy-performance trade-off for server farm management. PERFORMANCE 2010.

[QUESTA'10] V. Gupta, J. Dai, M. Harchol-Balter, and B. Zwart. On the inapproximability of M/G/K: why two moments of job size distribution are not enough. Queueing Systems, Vol 64, 2010.

[TR'08] V. Gupta, J. Dai, M. Harchol-Balter, and B. Zwart. The effect of higher moments of job size distribution on the performance of an M/G/K queueing system. Technical Report ,CMU, 2008.

[Sigmetrics'09] V. Gupta and M. Harchol-Balter. Self-adaptive admission control policies for resource-sharing systems. SIGMETRICS 2009.

[QUESTA'11] V. Gupta and T. Osogami, On Markov-Krein characterization of mean sojourn time in M/G/K.

Other Work

Time-varying systems	[Sigmetrics'06]	V. Gupta, M. Harchol-Balter, A. Scheller-Wolf, and U. Yechiali. Fundamental characteristics of queues with fluctuating load. Sigmetrics 2006
	[MAMA'08a]	V. Gupta, and P. Harrison. Fluid level in a reservoir with ON-OFF source. MAMA 2008.
Single Server Scheduling	[MAMA'08b]	V. Gupta. Finding the optimal quantum size: Sensitivity analysis of the M/G/1 round-robin queue. MAMA 2008.
	[Performance'10b]	V. Gupta, M. Burroughs, and M. Harchol-Balter. Analysis of scheduling policies under correlated job sizes. Performance 2010.
Distributed Data placement	[INFOCOM'10]	S. Borst, V. Gupta, and A. Walid. Distributed caching algorithms for Content Distribution Networks. INFOCOM 2010.
	[SOCC'10]	H. Amur, J. Cipar, V. Gupta, M. Kozuch, G.Ganger, and K. Schwan, Robust and flexible power-proportional storage. Symposium on Cloud Computing, 2010.
Epidemics	[INFOCOM'08]	M. Vojnovic, V. Gupta, T. Karagiannis, and C. Gkantsidis. Sampling strategies for epidemic-style information dissemination. INFOCOM 2008.
Stability analysis		A. Busic, V.Gupta, and J. Mairesse. Stability of the bipartite matching model. Under Submission